

# Rank Aggregation Using Scoring Rules

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# Rank Aggregation



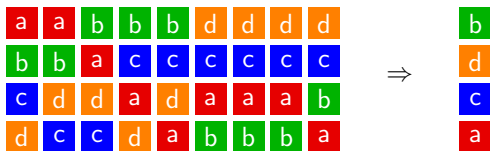
## Rank Aggregation

- ▶ Fundamental task in social choice theory.
- ▶ Numerous applications besides preference aggregation:  
meta-search engines, sports tournaments, multi-criteria decisions, ...

Best-known method: [Kemeny method](#)

- ▶ good axiomatic properties
- ▶ maximum likelihood estimator if there is a ground truth
- ▶ **but** hard to compute when there are many candidates
- ▶ and also hard to verify if a ranking is optimal (thus hard to audit)

# Rank Aggregation



## Project Goal

Understand the simple alternative: scoring-based rank aggregation rules.

not so simple alternative: approximation algorithms based on derandomization etc.

Scores: for example **plurality** (number of times ranked first), **Borda** (average rank), **veto** (number of times not ranked last)

- ▶ How to aggregate rankings based on scoring rules?
- ▶ What are the properties (axiomatic, computational efficiency) of scoring-based rank aggregation rules?
- ▶ How well do these rules perform in practice?

# Using Scoring Rules to Aggregate Rankings

Three ways of using scoring rules to aggregate rankings of  $m$  candidates:

## Score

Rank the candidates in order of their score.

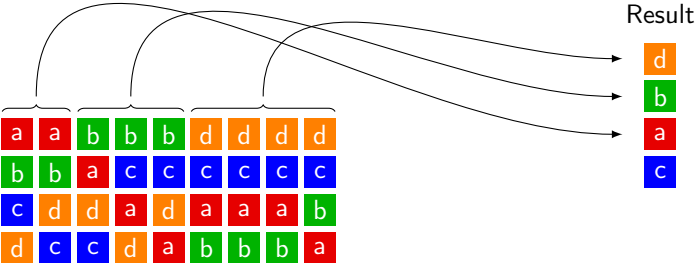
## Sequential-Winner

- ▶ Take the candidate  $c$  with highest score; rank it at position  $i := 1$ .
- ▶ **Repeat:** Delete  $c$  from profile; re-calculate scores; compute new candidate  $c$  with the highest score & rank it at position  $i := i + 1$ .

## Sequential-Loser

- ▶ Take the candidate  $c$  with lowest score; rank it at position  $i := m$ .
- ▶ **Repeat:** Delete  $c$  from profile; re-calculate scores; compute new candidate  $c$  with the lowest score & rank it at position  $i := i - 1$ .

# Examples: Plurality Score



# Examples: Sequential Plurality Winner

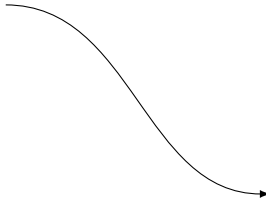
a	a	b	b	b	d	d	d	d
b	b	a	c	c	c	c	c	c
c	d	d	a	d	a	a	a	b
d	c	c	d	a	b	b	b	a

Result

- d
- c
- a
- b

# Examples: Sequential Plurality Winner

a	a	b	b	b	d	d	d	d
b	b	a	c	c	c	c	c	c
c	d	d	a	d	a	a	a	b
d	c	c	d	a	b	b	b	a



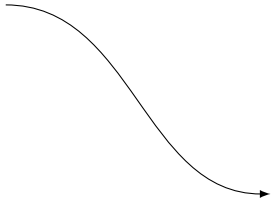
Result

- d
- c
- a
- b

# Examples: Sequential Plurality Winner

a	a	b	b	b	d	d	d	d
b	b	a	c	c	c	c	c	c
c	d	d	a	d	a	a	a	b
d	c	c	d	a	b	b	b	a

a	a	b	b	b				
b	b	a	c	c	c	c	c	c
c			a		a	a	a	b
	c	c		a	b	b	b	a



Result

d
c
a
b

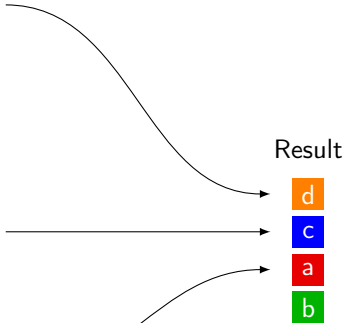


# Examples: Sequential Plurality Winner

a	a	b	b	b	d	d	d	d
b	b	a	c	c	c	c	c	c
c	d	d	a	d	a	a	a	b
d	c	c	d	a	b	b	b	a

a	a	b	b	b				
b	b	a	c	c	c	c	c	c
c			a		a	a	a	b
	c	c		a	b	b	b	a

a	a	b	b	b				
b	b	a						
			a		a	a	a	b
				a	b	b	b	a



# Examples: Sequential Plurality Loser

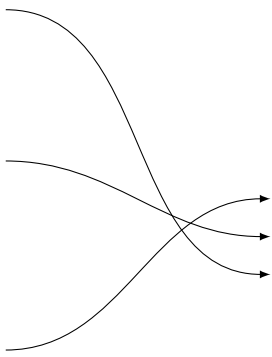
a	a	b	b	b	d	d	d	d
b	b	a	c	c	c	c	c	c
c	d	d	a	d	a	a	a	b
d	c	c	d	a	b	b	b	a

a	a	b	b	b	d	d	d	d
b	b	a						
	d	d	a	d	a	a	a	b
d			d	a	b	b	b	a

		b	b	b	d	d	d	d
b	b							
	d	d		d				b
d			d		b	b	b	

Result

b
d
a
c



# Axiomatic Properties

	Kemeny	Score			Seq.-Winner			Seq.-Loser		
		P	V	B	P	V	B	P	V	B
Independence at the top	✓				✓	✓	✓			
Independence at the bottom	✓							✓	✓	✓
Reinforcement	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓
Reinforcement at the top		✓	✓	✓	✓	✓	✓			
Reinforcement at the bottom		✓	✓	✓				✓	✓	✓
Condorcet winner at top	✓									✓
Copy majority	✓				✓				✓	
Independence of clones								✓		

K=Kemeny, P=Plurality, V=Veto, B=Borda

Much of this adapted from: Freeman, Brill, and Conitzer,  
*On the Axiomatic Characterization of Runoff Voting Rules*, AAAI 2014.

## Simulations I: How close to Kemeny?

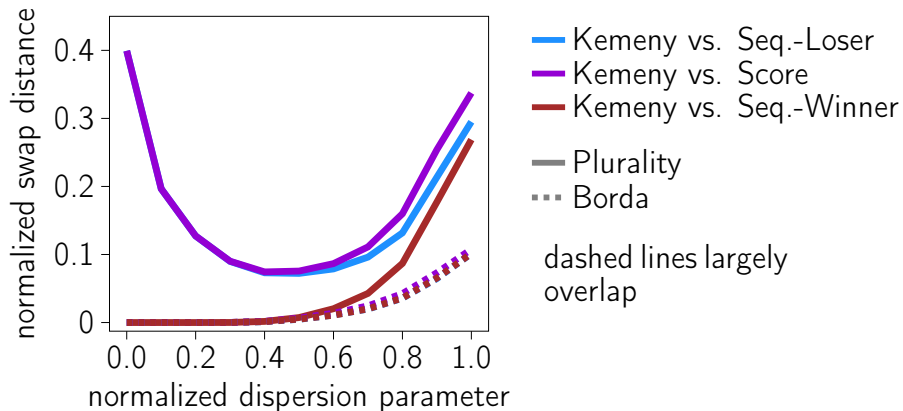


Figure: Pairwise average normalized swap distance on Mallows profiles with 10 candidates and 100 voters.

Seq.-Plurality-Winner produces rankings closest to Kemeny's ranking (among plurality-based rules). All Borda-based rules are quite close.

## Simulations II: How similar to each other?

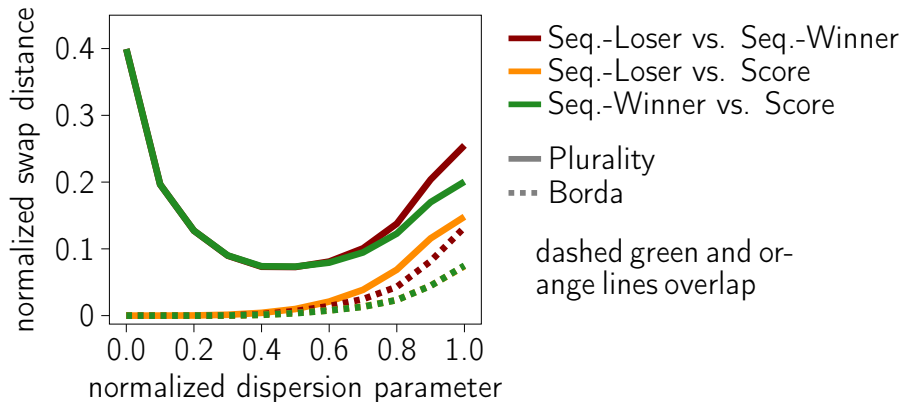
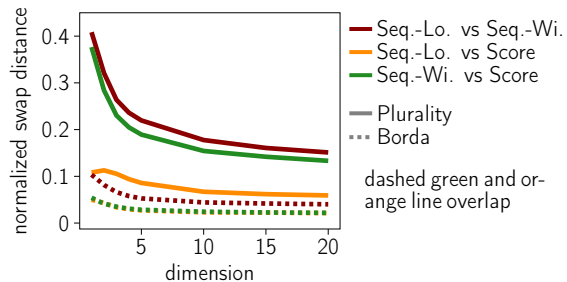


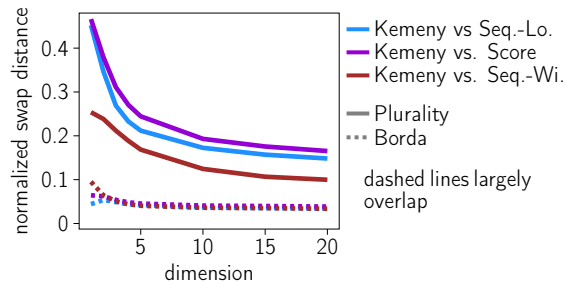
Figure: Pairwise average normalized swap distance on Mallows profiles with 10 candidates and 100 voters.

Seq.-Plurality-Loser and Plurality-Score are very similar, whereas Seq.-Plurality-Winner produces substantially different rankings.

## Simulations III: Euclidean profiles



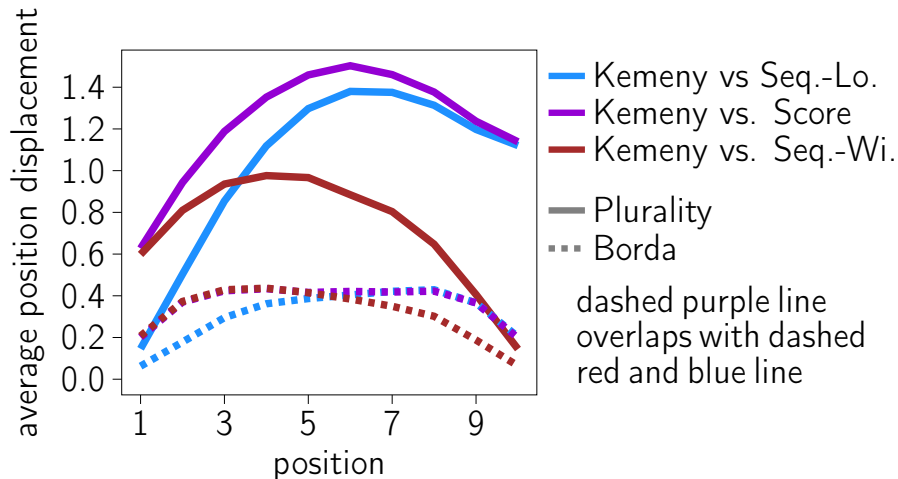
(a) Pairs of scoring-based methods



(b) Scoring-based vs. Kemeny ranking

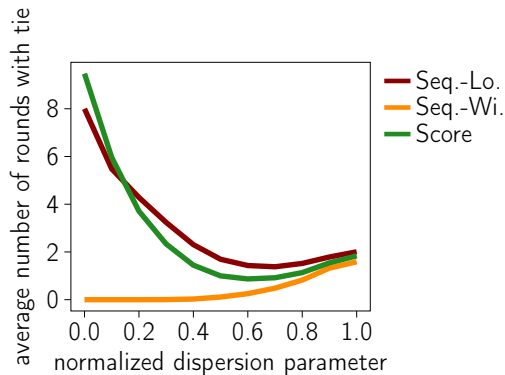
Figure: Pairwise average normalized swap distance on Euclidean profiles with 10 candidates and 100 voters.

## Simulations IV: Where are the disagreements?

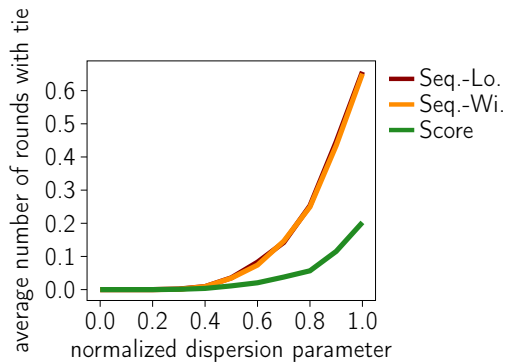


**Figure:** For pairs of rankings, average position displacement on each position for profiles generated using the Mallows model with  $\text{norm-}\phi = 0.8$  with 10 candidates and 100 voters.

## Simulations V: Frequency of ties



(a) Plurality



(b) Borda

**Figure:** Average number of rounds in which a tie occurs for Seq.-Plurality-Winner, Seq.-Plurality-Loser, and Plurality-Score on profiles with 10 candidates and 100 voters sampled from the Mallows model.



# Computational Complexity: Studied Problems

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## Position- $k$ Determination for social preference function $f$

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**Given:** A ranking profile  $P$  over candidate set  $C$ , a designated candidate  $d \in C$ , and an integer  $k \in [|C|]$ .

**Question:** Is there a ranking  $\succ$  selected by  $f$  on  $P$  where  $d$  is in position  $k$ , i.e.,  $\succ \in f(P)$  with  $\text{pos}(\succ, d) = k$ ?

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### Problem variants

- ▶ Top- $k$  Determination: Can a given candidate be placed in top- $k$  positions?
- ▶ Winner Determination: Special case for  $k = 1$ : Can a given candidate be placed in the top position?

## Computational Complexity: Sequential-Loser

Sequential-*-Loser		$n$	$m$
Plurality (STV)	NP-c.	FPT	FPT
Veto (Coombs)	NP-c.	W[1]-h., XP	FPT
Borda (Baldwin)	NP-c.	NP-c. ( $n = 8$ )	FPT

Table: Hardness: Winner Determination; Algorithms: Position- $k$  Determination.

- ▶ NP-hardness results stated or proven by Conitzer, Rognlie, and Xia (IJCAI '09) and Mattei, Narodytska, and Walsh (ECAI '14). Ours are a bit nicer and simpler IMHO.
- ▶ FPT:  $m!$  dependence is obvious, can get  $2^m$  by enumerating 'survival sets'.  
ETH lower bounds: cannot be solved in time  $2^{o(m)} \cdot \text{poly}(n, m)$ .
- ▶ Baldwin is hard for 8 voters: Use McGarvey's theorem very carefully. Same technique that shows Kemeny is hard for  $n = 4/7$  voters.

## Computational Complexity: Sequential Winner

Sequential-*Winner		$n$	$k$	$n + k$	$m$
Plurality	NP-c.	W[1]-h., XP	W[1]-h., XP	FPT	FPT
Veto	NP-c.	FPT	W[2]-h., XP	FPT	FPT
Borda	NP-c.	NP-h. ( $n = 8$ )	W[1]-h., XP	?	FPT

Table: Hardness: Top- $k$  Determination; Algorithms: Position- $k$  Determination