## Party Approval PAV Converges to Nash

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Let N be a set of n voters and C a set of m candidates (or parties). An approval profile is a collection  $A = (A_1, \ldots, A_n)$  of non-empty subsets  $A_i \subseteq C$  of approved parties for each voter  $i \in N$ .

Given a number of seats k, a seat allocation is a vector  $x = (x_1, \ldots, x_m)$  with  $x_i \in \mathbb{N}$  for each party  $i \in C$  such that  $\sum_{i=1}^{m} x_i = k$ . Given an approval profile A, and a number of seats k, the *PAV score* of a seat allocation x is defined as

$$PAV(A, x) = \sum_{i \in N} H(\sum_{j \in A_i} x_j),$$

where  $H(t) = \sum_{j=1}^{t} \frac{1}{j}$  is the harmonic number. The *PAV rule* selects a seat allocation x that maximizes the PAV score. Note that we are operating in the "party-approval" setting where each party can receive an arbitrary number of seats.

Let  $\Delta(C) = \{p \in [0,1]^m \mid \sum_{i=1}^m p_i = 1\}$  be the simplex of probability distributions over C. The Nash rule selects a distribution  $p \in \Delta(C)$  that maximizes the log Nash product

$$\operatorname{Nash}(A, p) = \sum_{i \in N} \log(\sum_{j \in A_i} p_j).$$

Note the following standard inequality connecting the harmonic and log functions:

$$\ln(t) + \gamma \le H(t) \le \ln(t) + \gamma + \frac{1}{2t},\tag{1}$$

where  $\gamma$  is the Euler-Mascheroni constant.

We can now show that the PAV rule converges to the Nash rule as  $k \to \infty$ .

**Theorem 1.** Let A be an approval profile. For each  $k \in \mathbb{N}$ , let  $x^k$  be the PAV allocation for A with k seats. Let  $p^k = \frac{x^k}{k}$  be the distribution over C induced by  $x^k$ . Since  $\Delta(C)$  is compact, the sequence  $p^k$  has a limit point  $p^{\infty} \in \Delta(C)$ . Let  $p^*$  be the Nash distribution for A. Then  $Nash(A, p^{\infty}) = Nash(A, p^*)$ .

*Proof.* Suppose for a contradiction that  $\operatorname{Nash}(A, p^{\infty}) < \operatorname{Nash}(A, p^*)$ . Then there exists a rational distribution q and some  $\varepsilon > 0$  such that  $\operatorname{Nash}(A, q) > \operatorname{Nash}(A, p^{\infty}) + \varepsilon$ .

Because the Nash objective is continuous and  $p^k \to p^{\infty}$ , we have  $\operatorname{Nash}(A, p^k) \to \operatorname{Nash}(A, p^{\infty})$ . Thus, for sufficiently large k, we have

$$Nash(A, p^k) < Nash(A, q) - \varepsilon/2.$$
<sup>(2)</sup>

Now, choose some  $k^*$  large enough such that

• (2) holds,

- $k^* > n^2/\varepsilon$ ,
- $k^*$  is a multiple of n, and
- $k^*$  is a multiple of the denominators of the rational numbers  $(q_j)_{j \in C}$ .

From the last property, there exists a seat allocation y such that  $q_j = y_j/k^*$  for all  $j \in C$ . We will show that y has a higher PAV score than  $x := x^{k^*}$ , a contradiction.

For each  $i \in N$ , write  $u_i = (\sum_{j \in A_i} x_j)/k^*$  for their utility under x as a fraction of seats, and  $u'_i = (\sum_{j \in A_i} y_j)/k^* = \sum_{j \in A_i} q_j$  for their utility under y and q. Recall that PAV satisfies EJR. Let  $i \in N$  be a voter. Note that  $S = \{i\}$  forms a group of size

Recall that PAV satisfies EJR. Let  $i \in N$  be a voter. Note that  $S = \{i\}$  forms a group of size  $|S| \geq \ell \cdot \frac{n}{k^*}$  with  $\ell = \lfloor \frac{k^*}{n} \rfloor$ , and S is obviously cohesive as approving at least one common party. Thus, EJR guarantees that

$$k^* u_i \ge \left\lfloor \frac{k^*}{n} \right\rfloor = \frac{k^*}{n},\tag{3}$$

where  $k^*u_i$  is the number of seats in x going to approved candidates, and where we can remove the floor because  $k^*$  is a multiple of n.

Now we have

$$PAV(A, x) = \sum_{i \in N} H(k^* u_i),$$
  
$$\leq \sum_{i \in N} \ln(k^* u_i) + \gamma + \frac{1}{2k^* u_i}$$
(by (1))

$$\leq \sum_{i \in N} \ln(k^* u_i) + \gamma + \frac{n}{2k^*} \tag{by (3)}$$

$$<\sum_{i\in N} \ln(k^*u_i) + \gamma + \frac{\varepsilon}{2n}$$
 (since  $k^* > n^2/\varepsilon$ )

$$= \operatorname{Nash}(A, p^{k}) + \sum_{i \in N} \ln(k) + \gamma + \frac{\varepsilon}{2n}$$
  
$$< \operatorname{Nash}(A, q) - \frac{\varepsilon}{2} + \sum_{i \in N} \ln(k^{*}) + \gamma + \frac{\varepsilon}{2n} \qquad (\text{from (2)})$$
  
$$< \operatorname{Nash}(A, q) - \frac{\varepsilon}{2} + \frac{\varepsilon}{2} + \sum_{i \in N} \ln(k^{*}) + \gamma$$
  
$$= \sum \ln(k^{*}u'_{i}) + \gamma$$

$$\leq \sum_{i \in N} H(k^* u'_i)$$

$$= PAV(A, y),$$
(by (1))

which is the desired contradiction.

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